

Formalizing Date Arithmetic and Statically Detecting Ambiguities for the Law

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Inria

 Université
de Lille

Some Catala News, and an Introduction

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Catala

- ▶ a DSL for computational laws

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- ▶ providing transparency
- ▶ easing maintenance
- ▶ through interdisciplinary work

News: Catala is getting closer to the real world

Partnerships with government agencies

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- ▶ Recognized as a priority program of support to public policies
- ▶ From 1 to 3 research engineers in the coming months (+ Denis)

Handling Date Computations in Catala

```
$ date -d "2024-01-31 + 1 month" +%F
```

```
$ date -d "2024-01-31 + 1 month" +%F  
2024-03-02
```

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$ date -d "2024-01-31 + 1 month" +%F  
2024-03-02  
$ date -d "2024-02-01 + 1 month" +%F
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```

Non-monotonic behavior?!

Different legal bodies and choices

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A wide variety of date semantics

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- ▶ 1 month = 30 days (Council of European Communities)
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 - 1 March in France, Germany, Hong-Kong
- ⇒ Formal, flexible semantics required!

Different legal bodies and choices

- ▶ 1 month = 30 days (Council of European Communities)
 - ▶ When do leapers become adults?
 - 28 February in New Zealand, Taiwan
 - 1 March in France, Germany, Hong-Kong
- ⇒ Formal, flexible semantics required! Focus on Gregorian calendar.

Outline

- 1 Semantics
- 2 Formalized Properties
- 3 Rounding-insensitivity Static Analysis
- 4 Case Study: French Housing Benefits
- 5 Conclusion

Semantics

values $v ::= (y, m, d) \mid \perp$
date unit $\delta ::= y \mid m \mid d$
expressions $e ::= v \mid e +_{\delta} n$

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date unit $\delta ::= y \mid m \mid d$
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$$\text{nb_days}(y, m) = \begin{cases} 29 & \text{if } m = 2 \wedge \text{is_leap}(y) \\ 28 & \text{if } m = 2 \wedge \neg \text{is_leap}(y) \\ 30 & \text{if } m \in \{ \mathbf{Apr}, \mathbf{Jun}, \mathbf{Sep}, \mathbf{Nov} \} \\ 31 & \text{otherwise} \end{cases}$$

Day additions with invalid day number propagate errors

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$$\frac{\text{ADD-DAYS-ERR1} \quad \text{day} < 1}{(y, m, \text{day}) +_d n \rightarrow \perp}$$

Day additions with invalid day number propagate errors

ADD-DAYS-ERR1

$day < 1$

$(y, m, day) +_d n \rightarrow \perp$

ADD-DAYS-ERR2

$day > nb_days(y, m)$

$(y, m, day) +_d n \rightarrow \perp$

Semantics – some cases of month addition

ADD-MONTH

$$1 \leq m + n \leq 12$$

$$(y, m, d) +_m n \rightarrow (y, m + n, d)$$

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ADD-MONTH-OVER

$$m + n > 12$$

$$(y, m, d) +_m n \rightarrow (y + 1, m, d) +_m (n - 12)$$

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$$(y, m, d) +_m n \rightarrow (y + 1, m, d) +_m (n - 12)$$

Similar cases for ADD-MONTH-UNDER, year, day addition.

$(2024, 01, 31) +_m 1 \rightarrow (2024, 02, 31)$

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Rounding to valid dates required!

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rounding mode $r ::= \uparrow | \downarrow | \perp$

expressions $e ::= v | e +_{\delta} n | \text{rnd}_r e$

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$\text{rnd}_\uparrow(2024, 02, 31) = (2024, 03, 01)$

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$\text{rnd}_\perp(2024, 02, 31) = \perp$

Coreutils-like rounding not defined here

ROUND-NOOP

$$\frac{1 \leq d \leq \text{nb_days}(y, m)}{\text{rnd}_r(y, m, d) \rightarrow (y, m, d)}$$

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ROUND-DOWN

$$d > \text{nb_days}(y, m)$$

$$\text{rnd}_\downarrow(y, m, d) \rightarrow (y, m, \text{nb_days}(y, m))$$

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ROUND-UP

$$\frac{d > \text{nb_days}(y, m) \quad (y, m, d) +_m 1 \xrightarrow{*} (y', m', d')}{\text{rnd}_\uparrow(y, m, d) \rightarrow (y', m', 1)}$$

ROUND-NOOP

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ROUND-ERR2

$$\frac{d > \text{nb_days}(y, m)}{\text{rnd}_\perp(y, m, d) \rightarrow \perp}$$

Date-period addition

Given a period (ys, ms, ds) :

$$e +_r (ys, ms, ds) ::= \text{rnd}_r((e +_y ys) +_m ms) +_d ds$$

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iff roundings e yield different values

Formalized Properties

Commutativity of addition

$$(2024, 03, 31) +_{\uparrow} 1m +_{\uparrow} 1d = (2024, 05, 01) +_{\uparrow} 1d = (2024, 05, 02)$$

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“Associativity” of addition

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$$(2024, 03, 31) +_r 2m = (2024, 05, 31)$$

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Well-formedness

For any date d , any period p , any value v , and $r \in \{\downarrow, \uparrow\}$, we have:

$$\text{valid}(d) \wedge d +_r p \xrightarrow{*} v \Rightarrow \text{valid}(v)$$

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Date addition is monotonic

For any dates d_1, d_2 , period p , $r \in \{\downarrow, \uparrow\}$, if $d_1 < d_2$, then $d_1 +_r p \leq d_2 +_r p$

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Loose bound in conclusion of monotonicity

$$(2024, 03, 30) +_{\downarrow} 1m = (2024, 04, 30) = (2024, 03, 31) +_{\downarrow} 1m$$

Rounding is monotonic

For all date d , period p :

1 $d +_{\downarrow} p \leq d +_{\uparrow} p$

2 $d +_{\perp} p \neq \perp \Rightarrow d +_{\downarrow} p = d +_{\uparrow} p = d +_{\perp} p$

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Equivalence of year and month addition

For all date d , for all integer n , $d +_y n = d +_m (12 * n)$.

Ambiguous month addition

For all valid date d , integer n such that $d +_m n \xrightarrow{*} (y, m, day)$:

$$\text{nb_days}(y, m) < \text{day} \Leftrightarrow \text{rnd}_{\perp}((y, m, \text{day})) \xrightarrow{*} \perp$$

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\Rightarrow core result needed for our static analysis

Rounding-insensitivity Static Analysis

Meaningful ambiguities

When rounding up or down doesn't change a computation

```
d + 1 month <= April 15 2024
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To reduce the need for costly legal interpretations

Rounding-insensitivity Static Analysis

Abstracting dates in a fixed rounding mode

- ▶ Defines addition, accessors, projection, lexicographic comparison

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$\mathbf{d}(d_1) \in [1, 31] \wedge \mathbf{m}(d_1) \in [1, 12] \wedge \mathbf{y}(d_1) = 2024$: all valid dates of 2024

Goal

Given a rounding mode, compute resulting dates from $d^\# +_m^\# n$, where $d^\#$ represents a set of dates.

Soundly derived from the ambiguous addition theorem.

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Algorithm: compute resulting month, year, then 4 cases:

- ▶ No rounding,
- ▶ Rounding, 30-day month,
- ▶ Rounding, non-leap years 28 Feb,
- ▶ Rounding, leap years, 29 Feb.

Partitioning used in practice.

YMD domain – month addition (II)

```
1 type case = expr * state
2 type cases = case list
3 let switch abs =
4   List.map (fun (cond : expr, k : state -> case) -> k (assume cond abs))
5
6 let add_months (r: rnd) ((d, m, y): var^3) (nb_m: int) (abs: state): cases =
7   let res_m: expr = 1 + (m - 1 + nb_m) % 12 in
8   let res_y: expr = y + (m - 1 + nb_m) / 12 in
9   switch abs
10  [
11    d > 30 && is_one_of res_m [Apr;Jun;Sep;Nov],
12    round r 30 res_m res_y;
13    d > 28 && res_m = Feb && not (is_leap res_y),
14    round r 28 res_m res_y;
15    d > 29 && res_m = Feb && is_leap res_y,
16    round r 29 res_m res_y;
17    mk_true,
18    mk_date d res_m res_y
19  ]
```

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date d1 = rand_date(); date d2 = d1 + 1 month; rounding down.
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¹Actually, $12y(d1) + m(d1) \leq 12y(d2) + 11 \wedge 12y(d2) \leq 12y(d1) + m(d1) + 1$

Rounding-insensitivity Static Analysis

Lifting to both rounding modes

Back to rounding-insensitivity detection

- ▶ Semantics on product programs with both rounding modes.

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- ▶ Inspired by Delmas, Ouadjaout, and Miné. “Static Analysis of Endian Portability by Abstract Interpretation”. SAS 2021.

Shallow variable duplication depending on their rounding mode.

Abstract double semantics

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date d1 = rand_date(); date d2 = d1 + 1 month; double semantics
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- ▶ 30-day month

$$d(d1) = 31, m(d1) \in \{ \text{Mar, May, Aug, Sep} \}$$

$$\downarrow d(d2) = 30, \downarrow m(d2) \in \{ \text{Apr, Jun, Sep, Nov} \}, \downarrow m(d2) = m(d1) + 1$$

$$\uparrow d(d2) = 1, \uparrow m(d2) \in \{ \text{May, Jul, Oct, Dec} \}, \uparrow m(d2) = m(d1) + 2$$

$$\downarrow y(d2) = \uparrow y(d2) = y(d1)$$



- ▶ Open-source static analysis platform



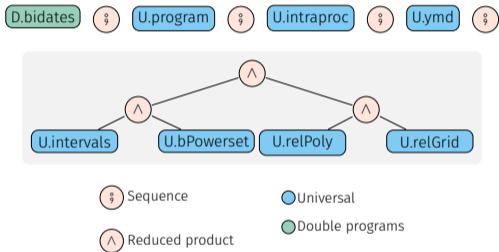
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Extracted sample from French housing benefits

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1 date current = rand_date();
2 date birthday = rand_date();
3 date intermediate = birthday + [2 years, 0 months, 0 days];
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```
5: assert(sync(current < limit));
      ^^^^^^^^^^^^^^^^^^^
```

Desynchronization detected: (current < limit). Hints:

```
↑month(limit) = 3, ↑day(limit) = 1, ↓month(limit) = 2, ↓day(limit) = 1,
↑month(intermediate) = 3, ↑day(intermediate) = 1, ↓month(intermediate) = 2,
↓day(intermediate) = 28, month(birthday) = 2, day(birthday) = 29,
year(birthday) = [4] 0, month(current) = 2, day(current) = [1,29],
year(current) = ↑year(intermediate) = ↑year(limit)
= ↓year(intermediate) = ↓year(limit) = year(birthday) + 2
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Computed, actual counter-example

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► current is in Feb. of year y

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- ▶ current is in Feb. of year y
- ▶ birthday is 29 Feb. of leap year y - 2
- ▶ intermediate is either 28 Feb. or 1 March of y

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Computed, actual counter-example

- ▶ `current` is in Feb. of year `y`
- ▶ `birthday` is 29 Feb. of leap year `y - 2`
- ▶ `intermediate` is either 28 Feb. or 1 March of `y`
- ▶ `limit` is either 1 Feb. or 1 March of `y`

Case Study: French Housing Benefits

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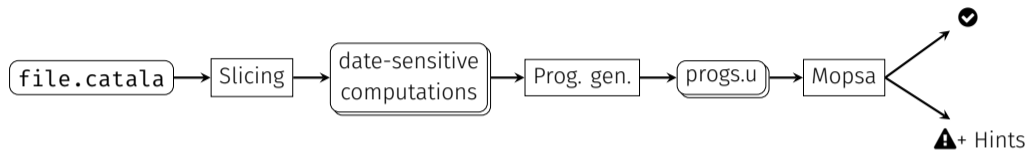
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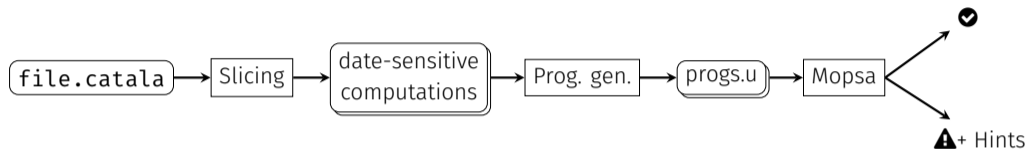
French Housing Benefits

20,000 Loc of Catala code (including text spec.)

Date ambiguity detection pipeline

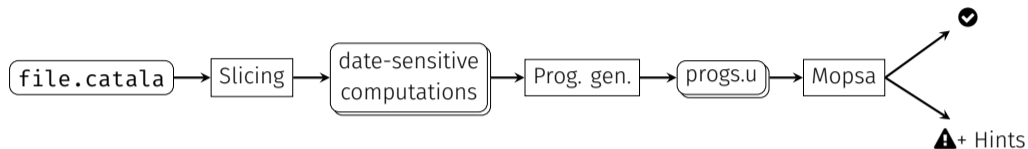


Date ambiguity detection pipeline



2 rounding-sensitive cases detected

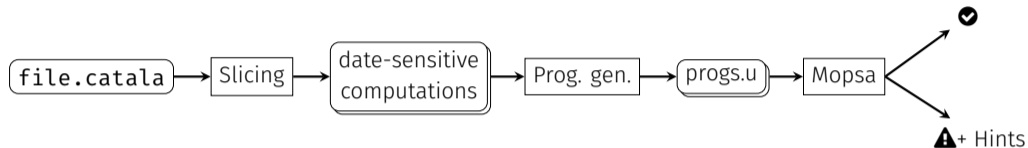
Date ambiguity detection pipeline



2 rounding-sensitive cases detected

Intra-scope extraction for now

Date ambiguity detection pipeline



2 rounding-sensitive cases detected

Intra-scope extraction for now

Manual inter-scope extraction

16 additional cases:

- ▶ 10 can be proved safe
(assuming `current_date` \geq 2023)
- ▶ Other are real issues

Conclusion

Survey of implementations

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Floating-point arithmetic

- ▶ FP widely used & more complex!
- ▶ Different rounding modes
- ▶ No analysis of rounding-sensitivity?

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Artefact & paper available! https://rmonat.fr/publication/24_esop/

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