

# AENEAS

*from Rust Programs  
to Pure Lambda Calculus*

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# Our experience with large-scale verification (1)

- 2016-2020: Project Everest (HACL\*, EverCrypt, Signal\*)
  - Low\*, a shallow embedding of C in F\*
  - **No separation logic:** modifies-clause theory *à la* Dafny
  - Heavy usage of **SMT-patterns** to reason about aliasing
  - Mostly **crypto primitives**: few buffers, straightforward aliasing, few branchings, arithmetic proofs
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  - 300k LoCs in F\*, 100k LoCs of generated C + ASM
- 2021: Noise\*, a protocol compiler for 59+ protocols
  - **High-level API**: sessions, state machines, peer library (imperative maps), etc.
  - Memory safety, functional correctness, security proofs
  - 55k LoCs of F\*; 5k LoCs of generated C per protocol instantiation
  - Still in Low\* (!)
  - Planned to verify WireGuard VPN on top of Noise\*: out of reach with current techno

# Our experience with large-scale verification (2)

What we no longer want to do:

```
// Yes, we need this!!!
let step_plus_minus_eq (step : ℕ) : Lemma ((step + 1) - 1 = step) = ()
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#restart-solver
-#push-options "--z3rlimit 2000 --ifuel 0"
-let mk_dstate_p_create #idc ssi initialize initiator r dvp epriv epub pid =
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```
// Auxiliary functions
inline let mk_session_p_create_uses_e_no_alloc = ...
inline let mk_session_p_create_uses_e_no_alloc_memzero = ... // calls mk_session_p_create_uses_e_no_alloc
inline let mk_session_p_create_uses_e = ... // mk_session_p_create_uses_e_no_alloc_memzero

let mk_session_p_create = ... // calls mk_session_p_create_uses_e
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2. Lack of interactivity (too much automation)
3. Non-modular (intrinsic) proofs

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- In C:

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void f (int *x, int *y);
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Dangling pointers?  
Aliasing?

B.live h x /\  
B.live h y /\  
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Dangling pointers?  
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B.live h x /\  
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- In Rust:  


```
fn f(x : &mut i32, y : &mut i32);
```

‘x’ and ‘y’ are live and disjoint



# Using Rust's type system to simplify verification

- **Not a new idea:**
  - **Prusti** (Viper): automate the application of memory reasoning rules
  - RustHornBelt/**Creusot**: “purification” through prophecy variables
  - **Verus**: early work (frontend *à la* Dafny)  
⇒ all target user-friendly frontends, **intrinsic proofs** and **high-automation** through SMT solvers via a *logical encoding*

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- **Our project:**
  - **Extrinsic proofs**
  - **Modularity**
  - Automated solvers are great, but we need an escape hatch: **interaction with tactics**

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- **Our project:**
  - **Extrinsic proofs**
  - **Modularity**
  - Automated solvers are great, but we need an escape hatch: **interaction with tactics**  
⇒ translate safe Rust programs to *idiomatic* pure lambda calculus (*executable*)  
⇒ generate code for several back-ends (Coq, F\*, HOL4...)
  - Similar to: **Electrolysis** (Rust to Lean), **Heapster** (LLVM to pure spec in Coq)

# Pure Translation – Returning Unknown Borrows

```
fn choose<'a>(b : bool, x : &'a mut i32, y : &'a mut i32)
    -> &'a mut i32
{
    if b { x }
    else { y }
}
```

```
// Usage example:
let mut x = 0;
let mut y = 1;
let z = choose(true, &mut x, &mut y); // (i)
*z = 2; // (ii)
// Lifetime 'a ends here - (iii)
assert!(x == 2);
assert!(y == 1);
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let choose (b : bool) (x : i32) (y : i32) : i32 =
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```

```
let x = 0 in
let y = 1 in
let z = choose true x y in // (i)
let z = 2 in // (ii)
let (x, y) = ?? in // (iii)
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let choose_fwd (b : bool) (x : i32) (y : i32) : i32 =
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- ⇒ Translation by means of forward and *backward* functions
- ⇒ Type of those functions only depends on *signature* of Rust functions

# Pure Translation – Lists & Recursion

In Rust:

```
let ls : List<T> = ...;
let x : &mut i32 = get_nth_mut(&mut ls, n);
*x = 4;
// end the borrow in x
...
```

# Pure Translation – Lists & Recursion

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In F\*:

```
let x = get_nth_mut_fwd ls n in // read
let x = 4 in
// end the borrow in x
let ls = get_nth_mut_back ls n x in // update
...
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let x = 4 in
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...
```

⇒ Read then update: **idiomatic** translation

# Pure Translation – PoC (function pointers)

Our translation is *entirely* type-based (hence modular)

In Rust:

```
fn apply_function(f: for<'a> fn(bool, &'a mut i32, &'a mut i32) -> &'a mut i32) -> i32 {  
    let mut x = 0;  
    let mut y = 1;  
    let pz = f(true, &mut x, &mut y);  
    *pz = 2;  
    // 'a ends here  
    return x;  
}
```

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In F\*:

```
let apply_function (f_fwd : bool -> i32 -> i32 -> i32)
                  (f_back : bool -> i32 -> i32 -> i32 -> (i32 * i32)) : i32 =
    let x = 0 in
    let y = 1 in
    let pz = f_fwd true x y in
    let pz = 2 in
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    x
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    let (x, y) = f_back true x y pz in
    x
```

- Useful because Rust is very higher-order (trait system)
- Same idea applies to **external dependencies** and **opaque modules**

# How does that work?

## **What do we need?**

- Know borrow graph at each point of the program
- Precisely abstract function calls

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## **3 steps:**

- **Operational semantics** for borrows
- **Regions + symbolic values** to handle functions calls;  
allows to perform symbolic executions
- Derivation of a **pure translation from a symbolic execution**

# Operational Semantics for Borrows - Mutable

```
let mut x = 0;           // (i)

let mut px1 = &mut x;    // (ii)

*px1 = 1;               // (iii)

let mut px2 = &mut (*px1); // (iv)

assert!(x == 1);        // (v)
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// (i)
x -> (0 : u32)
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x -> (0 : u32)
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x   -> (0 : u32)
px1 -> ?
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// (i)
x -> (0 : u32)
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x -> ?
px1 -> mut_borrow .. (0 : u32)
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// (i)
x -> (0 : u32)
// (ii)
x -> mut_loan 10
px1 -> mut_borrow 10 (0 : u32)
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let mut px1 = &mut x;    // (ii)
*px1 = 1;               // (iii)
let mut px2 = &mut (*px1); // (iv)
assert!(x == 1);        // (v)
```

```
// (i)
x -> (0 : u32)
// (ii)
x   -> mut_loan 10
px1 -> mut_borrow 10 (0 : u32)
// (iii)
x   -> mut_loan 10
px1 -> mut_borrow 10 (1 : u32)
// (iv)
x   -> mut_loan 10
px1 -> mut_borrow 10 (1 : u32)
px2 -> ?
```

# Operational Semantics for Borrows - Mutable

```
let mut x = 0;           // (i)
let mut px1 = &mut x;    // (ii)
*px1 = 1;               // (iii)
let mut px2 = &mut (*px1); // (iv)
assert!(x == 1);        // (v)
```

```
// (i)
x -> (0 : u32)
// (ii)
x   -> mut_loan 10
px1 -> mut_borrow 10 (0 : u32)
// (iii)
x   -> mut_loan 10
px1 -> mut_borrow 10 (1 : u32)
// (iv)
x   -> mut_loan 10
px1 -> mut_borrow 10 (1 : u32)
px2 -> ?
```

# Operational Semantics for Borrows - Mutable

```
let mut x = 0;           // (i)
let mut px1 = &mut x;    // (ii)
*px1 = 1;               // (iii)
let mut px2 = &mut (*px1); // (iv)
assert!(x == 1);        // (v)
```

```
// (i)
x -> (0 : u32)
// (ii)
x -> mut_loan 10
px1 -> mut_borrow 10 (0 : u32)
// (iii)
x -> mut_loan 10
px1 -> mut_borrow 10 (1 : u32)
// (iv)
x -> mut_loan 10
px1 -> mut_borrow 10 (mut_loan 11)
px2 -> mut_borrow 11 (1 : u32)
```

# Operational Semantics for Borrows - Mutable

```
let mut x = 0;           // (i)
let mut px1 = &mut x;    // (ii)
*px1 = 1;               // (iii)
let mut px2 = &mut (*px1); // (iv)
assert!(x == 1);        // (v)
```

```
// (i)
x -> (0 : u32)
// (ii)
x   -> mut_loan 10
px1 -> mut_borrow 10 (0 : u32)
// (iii)
x   -> mut_loan 10
px1 -> mut_borrow 10 (1 : u32)
// (iv)
x   -> mut_loan 10
px1 -> mut_borrow 10 (mut_loan 11)
px2 -> mut_borrow 11 (1 : u32)
```

# Operational Semantics for Borrows - Mutable

```
let mut x = 0;           // (i)
let mut px1 = &mut x;    // (ii)
*px1 = 1;               // (iii)
let mut px2 = &mut (*px1); // (iv)
assert!(x == 1);        // (v)
```

```
// (i)
x -> (0 : u32)
// (ii)
x   -> mut_loan 10
px1 -> mut_borrow 10 (0 : u32)
// (iii)
x   -> mut_loan 10
px1 -> mut_borrow 10 (1 : u32)
// (iv)
x   -> mut_loan 10
px1 -> mut_borrow 10 (mut_loan 11)
px2 -> mut_borrow 11 (1 : u32)
// (v)
x   -> mut_loan 10
px1 -> mut_borrow 10 (mut_loan 11)
px2 -> mut_borrow 11 (1 : u32)
```

# Operational Semantics for Borrows - Mutable

```
let mut x = 0;           // (i)
let mut px1 = &mut x;    // (ii)
*px1 = 1;               // (iii)
let mut px2 = &mut (*px1); // (iv)
assert!(x == 1);        // (v)
```

```
// (i)
x -> (0 : u32)
// (ii)
x   -> mut_loan 10
px1 -> mut_borrow 10 (0 : u32)
// (iii)
x   -> mut_loan 10
px1 -> mut_borrow 10 (1 : u32)
// (iv)
x   -> mut_loan 10
px1 -> mut_borrow 10 (mut_loan 11)
px2 -> mut_borrow 11 (1 : u32)
// (v)
x   -> mut_loan 10
px1 -> mut_borrow 10 (mut_loan 11)
px2 -> mut_borrow 11 (1 : u32)
```

# Operational Semantics for Borrows - Mutable

```
let mut x = 0;           // (i)
let mut px1 = &mut x;    // (ii)
*px1 = 1;               // (iii)
let mut px2 = &mut (*px1); // (iv)
assert!(x == 1);        // (v)
```

```
// (i)
x -> (0 : u32)
// (ii)
x   -> mut_loan 10
px1 -> mut_borrow 10 (0 : u32)
// (iii)
x   -> mut_loan 10
px1 -> mut_borrow 10 (1 : u32)
// (iv)
x   -> mut_loan 10
px1 -> mut_borrow 10 (mut_loan 11)
px2 -> mut_borrow 11 (1 : u32)
// (v)
x   -> mut_loan 10
px1 -> mut_borrow 10 (mut_loan 11)
px2 -> mut_borrow 11 (1 : u32)
```

# Operational Semantics for Borrows - Mutable

```
let mut x = 0;           // (i)
let mut px1 = &mut x;    // (ii)
*px1 = 1;               // (iii)
let mut px2 = &mut (*px1); // (iv)
assert!(x == 1);        // (v)
```

```
// (i)
x -> (0 : u32)
// (ii)
x   -> mut_loan 10
px1 -> mut_borrow 10 (0 : u32)
// (iii)
x   -> mut_loan 10
px1 -> mut_borrow 10 (1 : u32)
// (iv)
x   -> mut_loan 10
px1 -> mut_borrow 10 (mut_loan 11)
px2 -> mut_borrow 11 (1 : u32)
// (v)
x   -> mut_loan 10
px1 -> mut_borrow 10 (mut_loan 11)
px2 -> mut_borrow 11 (1 : u32)
```

# Operational Semantics for Borrows - Mutable

```
let mut x = 0;           // (i)
let mut px1 = &mut x;    // (ii)
*px1 = 1;               // (iii)
let mut px2 = &mut (*px1); // (iv)
assert!(x == 1);        // (v)
```

```
// (i)
x -> (0 : u32)
// (ii)
x   -> mut_loan 10
px1 -> mut_borrow 10 (0 : u32)
// (iii)
x   -> mut_loan 10
px1 -> mut_borrow 10 (1 : u32)
// (iv)
x   -> mut_loan 10
px1 -> mut_borrow 10 (mut_loan 11)
px2 -> mut_borrow 11 (1 : u32)
// (v)
x   -> mut_loan 10
px1 -> mut_borrow 10 (mut_loan 11)
px2 -> mut_borrow 11 (1 : u32)
```

# Operational Semantics for Borrows - Mutable

```
let mut x = 0;           // (i)  
  
let mut px1 = &mut x;    // (ii)  
  
*px1 = 1;              // (iii)  
  
let mut px2 = &mut (*px1); // (iv)  
  
assert!(x == 1);        // (v)
```

```
// (i)  
x -> (0 : u32)  
  
// (ii)  
x -> mut_loan 10  
px1 -> mut_borrow 10 (0 : u32)  
  
// (iii)  
x -> mut_loan 10  
px1 -> mut_borrow 10 (1 : u32)  
  
// (iv)  
x -> mut_loan 10  
px1 -> mut_borrow 10 (mut_loan 11)  
px2 -> mut_borrow 11 (1 : u32)  
  
// (v)  
x -> mut_loan 10  
px1 -> mut_borrow 10 (mut_loan 11)  
px2 -> mut_borrow 11 (1 : u32)
```

# Operational Semantics for Borrows - Mutable

```
let mut x = 0;           // (i)
let mut px1 = &mut x;    // (ii)
*px1 = 1;               // (iii)
let mut px2 = &mut (*px1); // (iv)
assert!(x == 1);        // (v)
```

```
// (i)
x -> (0 : u32)
// (ii)
x   -> mut_loan 10
px1 -> mut_borrow 10 (0 : u32)
// (iii)
x   -> mut_loan 10
px1 -> mut_borrow 10 (1 : u32)
// (iv)
x   -> mut_loan 10
px1 -> mut_borrow 10 (mut_loan 11)
px2 -> mut_borrow 11 (1 : u32)
// (v)
x   -> mut_loan 10
px1 -> mut_borrow 10 (mut_loan 11)
px2 -> mut_borrow 11 (1 : u32)
```

# Operational Semantics for Borrows - Mutable

```
let mut x = 0;           // (i)  
  
let mut px1 = &mut x;    // (ii)  
  
*px1 = 1;              // (iii)  
  
let mut px2 = &mut (*px1); // (iv)  
  
assert!(x == 1);        // (v)
```

```
// (i)  
x -> (0 : u32)  
  
// (ii)  
x -> mut_loan 10  
px1 -> mut_borrow 10 (0 : u32)  
  
// (iii)  
x -> mut_loan 10  
px1 -> mut_borrow 10 (1 : u32)  
  
// (iv)  
x -> mut_loan 10  
px1 -> mut_borrow 10 (mut_loan 11)  
px2 -> mut_borrow 11 (1 : u32)  
  
// (v)  
x -> mut_loan 10  
px1 -> mut_borrow 10 (1 : u32)  
px2 -> 1
```

# Operational Semantics for Borrows - Mutable

```
let mut x = 0;           // (i)
let mut px1 = &mut x;    // (ii)
*px1 = 1;               // (iii)
let mut px2 = &mut (*px1); // (iv)
assert!(x == 1);        // (v)
```

```
// (i)
x -> (0 : u32)
// (ii)
x   -> mut_loan 10
px1 -> mut_borrow 10 (0 : u32)
// (iii)
x   -> mut_loan 10
px1 -> mut_borrow 10 (1 : u32)
// (iv)
x   -> mut_loan 10
px1 -> mut_borrow 10 (mut_loan 11)
px2 -> mut_borrow 11 (1 : u32)
// (v)
x   -> mut_loan 10
px1 -> mut_borrow 10 (1 : u32)
px2 -> ⊥
```

# Operational Semantics for Borrows - Mutable

```
let mut x = 0;           // (i)
let mut px1 = &mut x;    // (ii)
*px1 = 1;               // (iii)
let mut px2 = &mut (*px1); // (iv)
assert!(x == 1);        // (v)
```

```
// (i)
x -> (0 : u32)
// (ii)
x   -> mut_loan 10
px1 -> mut_borrow 10 (0 : u32)
// (iii)
x   -> mut_loan 10
px1 -> mut_borrow 10 (1 : u32)
// (iv)
x   -> mut_loan 10
px1 -> mut_borrow 10 (mut_loan 11)
px2 -> mut_borrow 11 (1 : u32)
// (v)
x   -> (1 : u32)
px1 -> ⊥
px2 -> ⊥
```

# Operational Semantics for Borrows - Mutable

```
let mut x = 0;           // (i)
let mut px1 = &mut x;    // (ii)
*px1 = 1;               // (iii)
let mut px2 = &mut (*px1); // (iv)
assert!(x == 1);        // (v)
```

```
// (i)
x -> (0 : u32)
// (ii)
x   -> mut_loan 10
px1 -> mut_borrow 10 (0 : u32)
// (iii)
x   -> mut_loan 10
px1 -> mut_borrow 10 (1 : u32)
// (iv)
x   -> mut_loan 10
px1 -> mut_borrow 10 (mut_loan 11)
px2 -> mut_borrow 11 (1 : u32)
// (v)
x   -> (1 : u32)
px1 -> ⊥
px2 -> ⊥
```

- **Lazy** semantics of borrows
- Semantic analysis of borrows (Rust borrow checker is syntactic)

# Operational Semantics for Borrows - Shared

```
let x = 0;      // (i)  
  
let px1 = &x;    // (ii)  
  
let px2 = &x;    // (iii)
```

# Operational Semantics for Borrows - Shared

```
let x = 0;      // (i)  
  
let px1 = &x;    // (ii)  
  
let px2 = &x;    // (iii)
```

```
// (i)  
x -> (0 : i32)
```

# Operational Semantics for Borrows - Shared

```
let x = 0;      // (i)  
  
let px1 = &x;    // (ii)  
  
let px2 = &x;    // (iii)
```

```
// (i)  
x -> (0 : i32)  
  
// (ii)  
x  -> shared_loan {10} (0 : i32)  
px1 -> shared_borrow 10
```

# Operational Semantics for Borrows - Shared

```
let x = 0;          // (i)  
  
let px1 = &x;    // (ii)  
  
let px2 = &x;    // (iii)
```

```
// (i)  
x -> (0 : i32)  
  
// (ii)  
x   -> shared_loan {l0} (0 : i32)  
px1 -> shared_borrow l0  
  
// (iii)  
x   -> shared_loan {l0, l1} (0 : i32)  
px1 -> shared_borrow l0  
px2 -> shared_borrow l1
```

# Abstracting function calls

```
fn choose<'a>(b : bool, x : &'a mut i32, y : &'a mut i32) -> &'a mut i32
```

```
// Code
let mut x = 0;
let mut y = 1;
let px = &mut x;
let py = &mut y;
let z = choose(true, move px, move py);
*z = *z + 2;
assert!(x == 2);
assert!(y == 1);
```

```
// Env
```

# Abstracting function calls

```
fn choose<'a>(b : bool, x : &'a mut i32, y : &'a mut i32) -> &'a mut i32
```

```
// Code
let mut x = 0;
let mut y = 1;
let px = &mut x;
let py = &mut y;
let z = choose(true, move px, move py);
*z = *z + 2;
assert!(x == 2);
assert!(y == 1);
```

```
// Env
x -> mut_loan 10
y -> mut_loan 11
px -> mut_borrow 10 (0 : i32)
py -> mut_borrow 11 (1 : i32)
```

# Abstracting function calls

```
fn choose<'a>(b : bool, x : &'a mut i32, y : &'a mut i32) -> &'a mut i32
```

```
// Code
let mut x = 0;
let mut y = 1;
let px = &mut x;
let py = &mut y;
let z = choose(true, move px, move py);
*z = *z + 2;
assert!(x == 2);
assert!(y == 1);
```

```
// Env
x -> mut_loan 10
y -> mut_loan 11
px -> mut_borrow 10 (0 : i32)
py -> mut_borrow 11 (1 : i32)
```

# Abstracting function calls

```
fn choose<'a>(b : bool, x : &'a mut i32, y : &'a mut i32) -> &'a mut i32
```

```
// Code
let mut x = 0;
let mut y = 1;
let px = &mut x;
let py = &mut y;
let z = choose(true, move px, move py);
*z = *z + 2;
assert!(x == 2);
assert!(y == 1);
```

```
// Env
x -> mut_loan 10
y -> mut_loan 11
px -> mut_borrow 10 (0 : i32)
py -> mut_borrow 11 (1 : i32)
```

# Abstracting function calls

```
fn choose<'a>(b : bool, x : &'a mut i32, y : &'a mut i32) -> &'a mut i32
```

```
// Code
let mut x = 0;
let mut y = 1;
let px = &mut x;
let py = &mut y;
let z = choose(true, move px, move py);
*z = *z + 2;
assert!(x == 2);
assert!(y == 1);
```

```
// Env
x -> mut_loan 10
y -> mut_loan 11
px -> ⊥ // mut_borrow 10 (0 : i32)
py -> mut_borrow 11 (1 : i32)
```

# Abstracting function calls

```
fn choose<'a>(b : bool, x : &'a mut i32, y : &'a mut i32) -> &'a mut i32
```

```
// Code
let mut x = 0;
let mut y = 1;
let px = &mut x;
let py = &mut y;
let z = choose(true, move px, move py);
*z = *z + 2;
assert!(x == 2);
assert!(y == 1);
```

```
// Env
x -> mut_loan 10
y -> mut_loan 11
px -> ⊥ // mut_borrow 10 (0 : i32)
py -> ⊥ // mut_borrow 11 (1 : i32)
```

# Abstracting function calls

```
fn choose<'a>(b : bool, x : &'a mut i32, y : &'a mut i32) -> &'a mut i32
```

```
// Code
let mut x = 0;
let mut y = 1;
let px = &mut x;
let py = &mut y;
let z = choose(true, move px, move py);
*z = *z + 2;
assert!(x == 2);
assert!(y == 1);
```

```
// Env
x -> mut_loan 10
y -> mut_loan 11
px -> ⊥ // mut_borrow 10 (0 : i32)
py -> ⊥ // mut_borrow 11 (1 : i32)
z -> mut_borrow 12 (s0 : i32)
```

# Abstracting function calls

```
fn choose<'a>(b : bool, x : &'a mut i32, y : &'a mut i32) -> &'a mut i32
```

```
// Code
let mut x = 0;
let mut y = 1;
let px = &mut x;
let py = &mut y;
let z = choose(true, move px, move py);
*z = *z + 2;
assert!(x == 2);
assert!(y == 1);
```

```
// Env
x -> mut_loan 10
y -> mut_loan 11
px -> ⊥ // mut_borrow 10 (0 : i32)
py -> ⊥ // mut_borrow 11 (1 : i32)
z -> mut_borrow 12 (s0 : i32)
```

# Abstracting function calls

```
fn choose<'a>(b : bool, x : &'a mut i32, y : &'a mut i32) -> &'a mut i32
```

```
// Code
let mut x = 0;
let mut y = 1;
let px = &mut x;
let py = &mut y;
let z = choose(true, move px, move py);
*z = *z + 2;
assert!(x == 2);
assert!(y == 1);
```

```
// Env
x -> mut_loan 10
y -> mut_loan 11
px -> ⊥
py -> ⊥
z -> mut_borrow 12 (s0 : i32)
r0 {
    mut_borrow 10 (0 : i32)
    mut_borrow 11 (1 : i32)
    mut_loan 12
}
```

# Abstracting function calls

```
fn choose<'a>(b : bool, x : &'a mut i32, y : &'a mut i32) -> &'a mut i32
```

```
// Code
let mut x = 0;
let mut y = 1;
let px = &mut x;
let py = &mut y;
let z = choose(true, move px, move py);
*z = *z + 2;
assert!(x == 2);
assert!(y == 1);
```

```
// Env
x -> mut_loan 10
y -> mut_loan 11
px -> ⊥
py -> ⊥
z -> mut_borrow 12 (s0 : i32)
r0 {
    mut_borrow 10 (0 : i32)
    mut_borrow 11 (1 : i32)
    mut_loan 12
}
```

# Abstracting function calls

```
fn choose<'a>(b : bool, x : &'a mut i32, y : &'a mut i32) -> &'a mut i32
```

```
// Code
let mut x = 0;
let mut y = 1;
let px = &mut x;
let py = &mut y;
let z = choose(true, move px, move py);
*z = *z + 2;
assert!(x == 2);
assert!(y == 1);
```

```
// Env
x -> mut_loan 10
y -> mut_loan 11
px -> ⊥
py -> ⊥
z -> mut_borrow 12 (s0 : i32)
r0 {
    mut_borrow 10 (0 : i32)
    mut_borrow 11 (1 : i32)
    mut_loan 12
}
```

# Abstracting function calls

```
fn choose<'a>(b : bool, x : &'a mut i32, y : &'a mut i32) -> &'a mut i32
```

```
// Code
let mut x = 0;
let mut y = 1;
let px = &mut x;
let py = &mut y;
let z = choose(true, move px, move py);
*z = *z + 2;
assert!(x == 2);
assert!(y == 1);
```

```
// Env
x -> mut_loan 10
y -> mut_loan 11
px -> ⊥
py -> ⊥
z -> mut_borrow 12 (s0 : i32)
r0 {
    mut_borrow 10 (0 : i32)
    mut_borrow 11 (1 : i32)
    mut_loan 12
}
```

# Abstracting function calls

```
fn choose<'a>(b : bool, x : &'a mut i32, y : &'a mut i32) -> &'a mut i32
```

```
// Code
let mut x = 0;
let mut y = 1;
let px = &mut x;
let py = &mut y;
let z = choose(true, move px, move py);
*z = *z + 2;
assert!(x == 2);
assert!(y == 1);
```

```
// Env
x -> mut_loan 10
y -> mut_loan 11
px -> ⊥
py -> ⊥
z -> mut_borrow 12 (s1 : i32)
r0 {
    mut_borrow 10 (0 : i32)
    mut_borrow 11 (1 : i32)
    mut_loan 12
}
```

# Abstracting function calls

```
fn choose<'a>(b : bool, x : &'a mut i32, y : &'a mut i32) -> &'a mut i32
```

```
// Code
let mut x = 0;
let mut y = 1;
let px = &mut x;
let py = &mut y;
let z = choose(true, move px, move py);
*z = *z + 2;
assert!(x == 2);
assert!(y == 1);
```

```
// Env
x -> mut_loan 10
y -> mut_loan 11
px -> ⊥
py -> ⊥
z -> mut_borrow 12 (s1 : i32)
r0 {
    mut_borrow 10 (0 : i32)
    mut_borrow 11 (1 : i32)
    mut_loan 12
}
```

# Abstracting function calls

```
fn choose<'a>(b : bool, x : &'a mut i32, y : &'a mut i32) -> &'a mut i32
```

```
// Code
let mut x = 0;
let mut y = 1;
let px = &mut x;
let py = &mut y;
let z = choose(true, move px, move py);
*z = *z + 2;
assert!(x == 2);
assert!(y == 1);
```

```
// Env
x -> mut_loan 10
y -> mut_loan 11
px -> ⊥
py -> ⊥
z -> mut_borrow 12 (s1 : i32)
r0 {
    mut_borrow 10 (0 : i32)
    mut_borrow 11 (1 : i32)
    mut_loan 12
}
```

# Abstracting function calls

```
fn choose<'a>(b : bool, x : &'a mut i32, y : &'a mut i32) -> &'a mut i32
```

```
// Code
let mut x = 0;
let mut y = 1;
let px = &mut x;
let py = &mut y;
let z = choose(true, move px, move py);
*z = *z + 2;
assert!(x == 2);
assert!(y == 1);
```

```
// Env
x -> mut_loan 10
y -> mut_loan 11
px -> ⊥
py -> ⊥
z -> mut_borrow 12 (s1 : i32)
r0 {
    mut_borrow 10 (0 : i32)
    mut_borrow 11 (1 : i32)
    mut_loan 12
}
```

# Abstracting function calls

```
fn choose<'a>(b : bool, x : &'a mut i32, y : &'a mut i32) -> &'a mut i32
```

```
// Code
let mut x = 0;
let mut y = 1;
let px = &mut x;
let py = &mut y;
let z = choose(true, move px, move py);
*z = *z + 2;
assert!(x == 2);
assert!(y == 1);
```

```
// Env
x -> mut_loan 10
y -> mut_loan 11
px -> ⊥
py -> ⊥
z -> ⊥
r0 {
    mut_borrow 10 (0 : i32)
    mut_borrow 11 (1 : i32)
    (s1 : i32)
}
```

# Abstracting function calls

```
fn choose<'a>(b : bool, x : &'a mut i32, y : &'a mut i32) -> &'a mut i32
```

```
// Code
let mut x = 0;
let mut y = 1;
let px = &mut x;
let py = &mut y;
let z = choose(true, move px, move py);
*z = *z + 2;
assert!(x == 2);
assert!(y == 1);
```

```
// Env
x -> mut_loan 10
y -> mut_loan 11
px -> ⊥
py -> ⊥
z -> ⊥
r0 {
    mut_borrow 10 (0 : i32)
    mut_borrow 11 (1 : i32)
    (s1 : i32)
}
```

# Abstracting function calls

```
fn choose<'a>(b : bool, x : &'a mut i32, y : &'a mut i32) -> &'a mut i32
```

```
// Code
let mut x = 0;
let mut y = 1;
let px = &mut x;
let py = &mut y;
let z = choose(true, move px, move py);
*z = *z + 2;
assert!(x == 2);
assert!(y == 1);
```

```
// Env
x -> mut_loan 10
y -> mut_loan 11
px -> ⊥
py -> ⊥
z -> ⊥
r0 {
    mut_borrow 10 (s2 : i32)
    mut_borrow 11 (s3 : i32)
    (s1 : i32)
}
```

# Abstracting function calls

```
fn choose<'a>(b : bool, x : &'a mut i32, y : &'a mut i32) -> &'a mut i32
```

```
// Code
let mut x = 0;
let mut y = 1;
let px = &mut x;
let py = &mut y;
let z = choose(true, move px, move py);
*z = *z + 2;
assert!(x == 2);
assert!(y == 1);
```

```
// Env
x -> (s2 : i32)
y -> (s3 : i32)
px -> ⊥
py -> ⊥
z -> ⊥
r0 {
    ⊥ // gave back: s2
    ⊥ // gave back: s3
    (s1 : i32)
}
```

# Abstracting function calls

```
fn choose<'a>(b : bool, x : &'a mut i32, y : &'a mut i32) -> &'a mut i32
```

```
// Code
let mut x = 0;
let mut y = 1;
let px = &mut x;
let py = &mut y;
let z = choose(true, move px, move py);
*z = *z + 2;
assert!(x == 2);
assert!(y == 1);
```

```
// Env
x -> (s2 : i32)
y -> (s3 : i32)
px -> ⊥
py -> ⊥
z -> ⊥
r0 {
    ⊥ // gave back: s2
    ⊥ // gave back: s3
    (s1 : i32)
}
```

# Abstracting function calls + Translation

```
fn choose<'a>(b : bool, x : &'a mut i32, y : &'a mut i32) -> &'a mut i32
```

```
// Code
let mut x = 0;
let mut y = 1;
let px = &mut x;
let py = &mut y;
let z = choose(true, move px, move py);
*z = *z + 2;
assert!(x == 2);
assert!(y == 1);
```

```
// Env
```

```
// Translation
[.]
```

# Abstracting function calls + Translation

```
fn choose<'a>(b : bool, x : &'a mut i32, y : &'a mut i32) -> &'a mut i32
```

```
// Code
let mut x = 0;
let mut y = 1;
let px = &mut x;
let py = &mut y;
let z = choose(true, move px, move py);
*z = *z + 2;
assert!(x == 2);
assert!(y == 1);
```

```
// Env
x -> mut_loan 10
y -> mut_loan 11
px -> mut_borrow 10 (0 : i32)
py -> mut_borrow 11 (1 : i32)
```

```
// Translation
[.]
```

# Abstracting function calls + Translation

```
fn choose<'a>(b : bool, x : &'a mut i32, y : &'a mut i32) -> &'a mut i32
```

```
// Code
let mut x = 0;
let mut y = 1;
let px = &mut x;
let py = &mut y;
let z = choose(true, move px, move py);
*z = *z + 2;
assert!(x == 2);
assert!(y == 1);
```

```
// Env
x -> mut_loan 10
y -> mut_loan 11
px -> mut_borrow 10 (0 : i32)
py -> mut_borrow 11 (1 : i32)
```

```
// Translation
... <-- choose_fwd .... . . . ;
[.]
```

# Abstracting function calls + Translation

```
fn choose<'a>(b : bool, x : &'a mut i32, y : &'a mut i32) -> &'a mut i32
```

```
// Code
let mut x = 0;
let mut y = 1;
let px = &mut x;
let py = &mut y;
let z = choose(true, move px, move py);
*z = *z + 2;
assert!(x == 2);
assert!(y == 1);
```

```
// Env
x -> mut_loan 10
y -> mut_loan 11
px -> mut_borrow 10 (0 : i32)
py -> mut_borrow 11 (1 : i32)
```

```
// Translation
.. <-- choose_fwd true . .;
[.]
```

# Abstracting function calls + Translation

```
fn choose<'a>(b : bool, x : &'a mut i32, y : &'a mut i32) -> &'a mut i32
```

```
// Code
let mut x = 0;
let mut y = 1;
let px = &mut x;
let py = &mut y;
let z = choose(true, move px, move py);
*z = *z + 2;
assert!(x == 2);
assert!(y == 1);
```

```
// Env
x -> mut_loan 10
y -> mut_loan 11
px -> ⊥ // mut_borrow 10 (0 : i32)
py -> mut_borrow 11 (1 : i32)
```

```
// Translation
.. <-- choose_fwd true 0 .;
[.]
```

# Abstracting function calls + Translation

```
fn choose<'a>(b : bool, x : &'a mut i32, y : &'a mut i32) -> &'a mut i32
```

```
// Code
let mut x = 0;
let mut y = 1;
let px = &mut x;
let py = &mut y;
let z = choose(true, move px, move py);
*z = *z + 2;
assert!(x == 2);
assert!(y == 1);
```

```
// Env
x -> mut_loan 10
y -> mut_loan 11
px -> ⊥ // mut_borrow 10 (0 : i32)
py -> ⊥ // mut_borrow 11 (1 : i32)
```

```
// Translation
.. <-- choose_fwd true 0 1;
[.]
```

# Abstracting function calls + Translation

```
fn choose<'a>(b : bool, x : &'a mut i32, y : &'a mut i32) -> &'a mut i32
```

```
// Code
let mut x = 0;
let mut y = 1;
let px = &mut x;
let py = &mut y;
let z = choose(true, move px, move py);
*z = *z + 2;
assert!(x == 2);
assert!(y == 1);
```

```
// Env
x -> mut_loan 10
y -> mut_loan 11
px -> ⊥ // mut_borrow 10 (0 : i32)
py -> ⊥ // mut_borrow 11 (1 : i32)
z -> mut_borrow 12 (s0 : i32)
```

```
// Translation
s0 <- choose_fwd true 0 1;
[.]
```

# Abstracting function calls + Translation

```
fn choose<'a>(b : bool, x : &'a mut i32, y : &'a mut i32) -> &'a mut i32
```

```
// Code
let mut x = 0;
let mut y = 1;
let px = &mut x;
let py = &mut y;
let z = choose(true, move px, move py);
*z = *z + 2;
assert!(x == 2);
assert!(y == 1);
```

```
// Env
x -> mut_loan 10
y -> mut_loan 11
px -> ⊥ // mut_borrow 10 (0 : i32)
py -> ⊥ // mut_borrow 11 (1 : i32)
z -> mut_borrow 12 (s0 : i32)
```

```
// Translation
s0 <-- choose_fwd true 0 1;
[.]
```

# Abstracting function calls + Translation

```
fn choose<'a>(b : bool, x : &'a mut i32, y : &'a mut i32) -> &'a mut i32
```

```
// Code
let mut x = 0;
let mut y = 1;
let px = &mut x;
let py = &mut y;
let z = choose(true, move px, move py);
*z = *z + 2;
assert!(x == 2);
assert!(y == 1);
```

```
// Env
x -> mut_loan 10
y -> mut_loan 11
px -> ⊥
py -> ⊥
z -> mut_borrow 12 (s0 : i32)
r0 {
    mut_borrow 10 (0 : i32)
    mut_borrow 11 (1 : i32)
    mut_loan 12
}
```

```
// Translation
s0 <-- choose_fwd true 0 1;
[.]
```

# Abstracting function calls + Translation

```
fn choose<'a>(b : bool, x : &'a mut i32, y : &'a mut i32) -> &'a mut i32
```

```
// Code
let mut x = 0;
let mut y = 1;
let px = &mut x;
let py = &mut y;
let z = choose(true, move px, move py);
*z = *z + 2;
assert!(x == 2);
assert!(y == 1);
```

```
// Env
x -> mut_loan 10
y -> mut_loan 11
px -> ⊥
py -> ⊥
z -> mut_borrow 12 (s0 : i32)
r0 {
    mut_borrow 10 (0 : i32)
    mut_borrow 11 (1 : i32)
    mut_loan 12
}
```

```
// Translation
s0 <- choose_fwd true 0 1;
... <- i32_add ... .;
[.]
```

# Abstracting function calls + Translation

```
fn choose<'a>(b : bool, x : &'a mut i32, y : &'a mut i32) -> &'a mut i32
```

```
// Code
let mut x = 0;
let mut y = 1;
let px = &mut x;
let py = &mut y;
let z = choose(true, move px, move py);
*z = *z + 2;
assert!(x == 2);
assert!(y == 1);
```

```
// Env
x -> mut_loan 10
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px -> ⊥
py -> ⊥
z -> mut_borrow 12 (s0 : i32)
r0 {
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    mut_borrow 11 (1 : i32)
    mut_loan 12
}
```

```
// Translation
s0 <- choose_fwd true 0 1;
.. <- i32_add s0 .;
[.]
```

# Abstracting function calls + Translation

```
fn choose<'a>(b : bool, x : &'a mut i32, y : &'a mut i32) -> &'a mut i32
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// Code
let mut x = 0;
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let py = &mut y;
let z = choose(true, move px, move py);
*z = *z + 2;
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```

```
// Env
x -> mut_loan 10
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```
// Translation
s0 <- choose_fwd true 0 1;
.. <- i32_add s0 2;
[.]
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# Abstracting function calls + Translation

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fn choose<'a>(b : bool, x : &'a mut i32, y : &'a mut i32) -> &'a mut i32
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let mut x = 0;
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let py = &mut y;
let z = choose(true, move px, move py);
*z = *z + 2;
assert!(x == 2);
assert!(y == 1);
```

```
// Env
x -> mut_loan 10
y -> mut_loan 11
px -> ⊥
py -> ⊥
z -> mut_borrow 12 (s1 : i32)
r0 {
    mut_borrow 10 (0 : i32)
    mut_borrow 11 (1 : i32)
    mut_loan 12
}
```

```
// Translation
s0 <- choose_fwd true 0 1;
s1 <- i32_add s0 2;
[.]
```

# Abstracting function calls + Translation

```
fn choose<'a>(b : bool, x : &'a mut i32, y : &'a mut i32) -> &'a mut i32
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let mut x = 0;
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let z = choose(true, move px, move py);
*z = *z + 2;
assert!(x == 2);
assert!(y == 1);
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// Env
x -> mut_loan 10
y -> mut_loan 11
px -> ⊥
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z -> mut_borrow 12 (s1 : i32)
r0 {
    mut_borrow 10 (0 : i32)
    mut_borrow 11 (1 : i32)
    mut_loan 12
}
```

```
// Translation
s0 <- choose_fwd true 0 1;
s1 <- i32_add s0 2;
[.]
```

# Abstracting function calls + Translation

```
fn choose<'a>(b : bool, x : &'a mut i32, y : &'a mut i32) -> &'a mut i32
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```
// Code
let mut x = 0;
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let px = &mut x;
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let z = choose(true, move px, move py);
*z = *z + 2;
assert!(x == 2);
assert!(y == 1);
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```
// Env
x -> mut_loan 10
y -> mut_loan 11
px -> ⊥
py -> ⊥
z -> mut_borrow 12 (s1 : i32)
r0 {
    mut_borrow 10 (0 : i32)
    mut_borrow 11 (1 : i32)
    mut_loan 12
}
```

```
// Translation
s0 <- choose_fwd true 0 1;
s1 <- i32_add s0 2;
..... <- choose_back true 0 1 ..;
[.]
```

# Abstracting function calls + Translation

```
fn choose<'a>(b : bool, x : &'a mut i32, y : &'a mut i32) -> &'a mut i32
```

```
// Code
let mut x = 0;
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let px = &mut x;
let py = &mut y;
let z = choose(true, move px, move py);
*z = *z + 2;
assert!(x == 2);
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```

```
// Env
x -> mut_loan 10
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px -> ⊥
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z -> mut_borrow 12 (s1 : i32)
r0 {
    mut_borrow 10 (0 : i32)
    mut_borrow 11 (1 : i32)
    mut_loan 12
}
```

```
// Translation
s0 <- choose_fwd true 0 1;
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..... <- choose_back true 0 1 ..;
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# Abstracting function calls + Translation

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fn choose<'a>(b : bool, x : &'a mut i32, y : &'a mut i32) -> &'a mut i32
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// Code
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*z = *z + 2;
assert!(x == 2);
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```

```
// Env
x -> mut_loan 10
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px -> ⊥
py -> ⊥
z -> ⊥
r0 {
    mut_borrow 10 (0 : i32)
    mut_borrow 11 (1 : i32)
    (s1 : i32)
}
```

```
// Translation
s0 <- choose_fwd true 0 1;
s1 <- i32_add s0 2;
..... <- choose_back true 0 1 s1;
[.]
```

# Abstracting function calls + Translation

```
fn choose<'a>(b : bool, x : &'a mut i32, y : &'a mut i32) -> &'a mut i32
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```
// Code
let mut x = 0;
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let px = &mut x;
let py = &mut y;
let z = choose(true, move px, move py);
*z = *z + 2;
assert!(x == 2);
assert!(y == 1);
```

```
// Env
x -> mut_loan 10
y -> mut_loan 11
px -> ⊥
py -> ⊥
z -> ⊥
r0 {
    mut_borrow 10 (0 : i32)
    mut_borrow 11 (1 : i32)
    (s1 : i32)
}
```

```
// Translation
s0 <- choose_fwd true 0 1;
s1 <- i32_add s0 2;
..... <- choose_back true 0 1 s1;
[.]
```

# Abstracting function calls + Translation

```
fn choose<'a>(b : bool, x : &'a mut i32, y : &'a mut i32) -> &'a mut i32
```

```
// Code
let mut x = 0;
let mut y = 1;
let px = &mut x;
let py = &mut y;
let z = choose(true, move px, move py);
*z = *z + 2;
assert!(x == 2);
assert!(y == 1);
```

```
// Env
x -> mut_loan 10
y -> mut_loan 11
px -> ⊥
py -> ⊥
z -> ⊥
r0 {
    mut_borrow 10 (s2 : i32)
    mut_borrow 11 (s3 : i32)
    (s1 : i32)
}
```

```
// Translation
s0 <- choose_fwd true 0 1;
s1 <- i32_add s0 2;
..... <- choose_back true 0 1 s1;
[.]
```

# Abstracting function calls + Translation

```
fn choose<'a>(b : bool, x : &'a mut i32, y : &'a mut i32) -> &'a mut i32
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```
// Code
let mut x = 0;
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let px = &mut x;
let py = &mut y;
let z = choose(true, move px, move py);
*z = *z + 2;
assert!(x == 2);
assert!(y == 1);
```

```
// Env
x -> (s2 : i32)
y -> (s3 : i32)
px -> ⊥
py -> ⊥
z -> ⊥
r0 {
    ⊥ // gave back: s2
    ⊥ // gave back: s3
    (s1 : i32)
}
```

```
// Translation
s0 <- choose_fwd true 0 1;
s1 <- i32_add s0 2;
(s2, s3) <- choose_back true 0 1 s1;
[.]
```

# Abstracting function calls + Translation

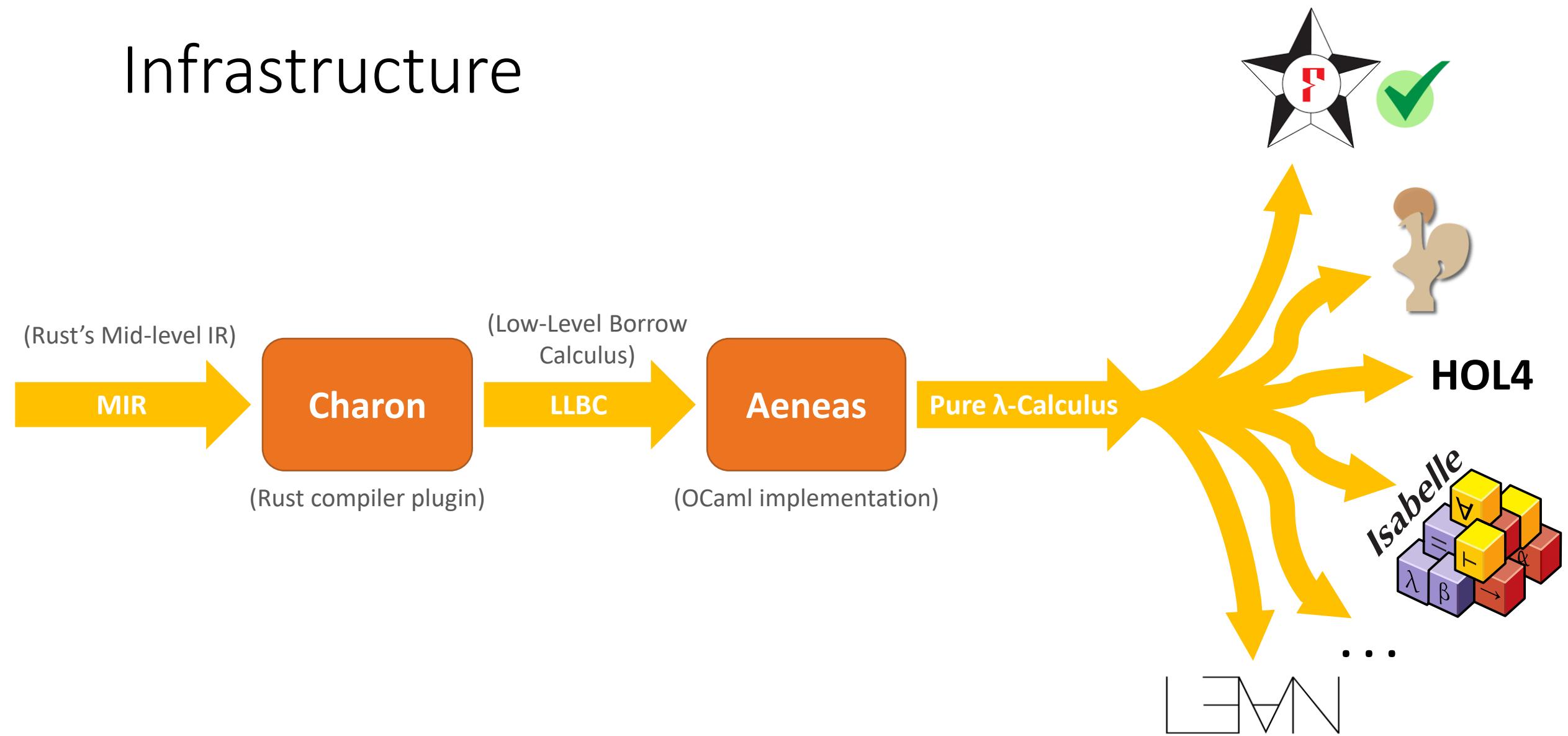
```
fn choose<'a>(b : bool, x : &'a mut i32, y : &'a mut i32) -> &'a mut i32
```

```
// Code
let mut x = 0;
let mut y = 1;
let px = &mut x;
let py = &mut y;
let z = choose(true, move px, move py);
*z = *z + 2;
assert!(x == 2);
assert!(y == 1);
```

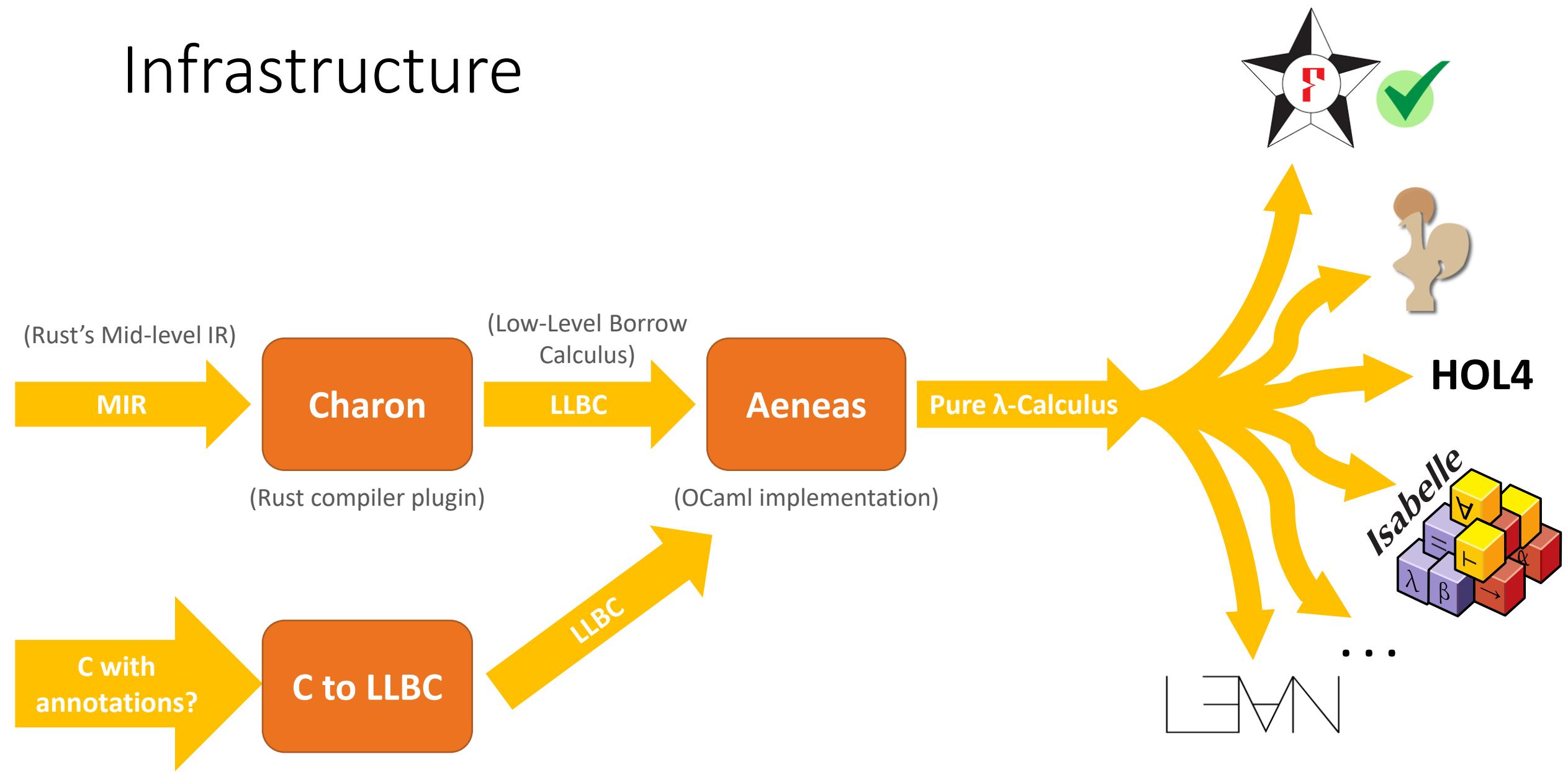
```
// Env
x -> (s2 : i32)
y -> (s3 : i32)
px -> ⊥
py -> ⊥
z -> ⊥
r0 {
    ⊥ // gave back: s2
    ⊥ // gave back: s3
    (s1 : i32)
}
```

```
// Translation
s0 <- choose_fwd true 0 1;
s1 <- i32_add s0 2;
(s2, s3) <- choose_back true 0 1 s1;
[.]
```

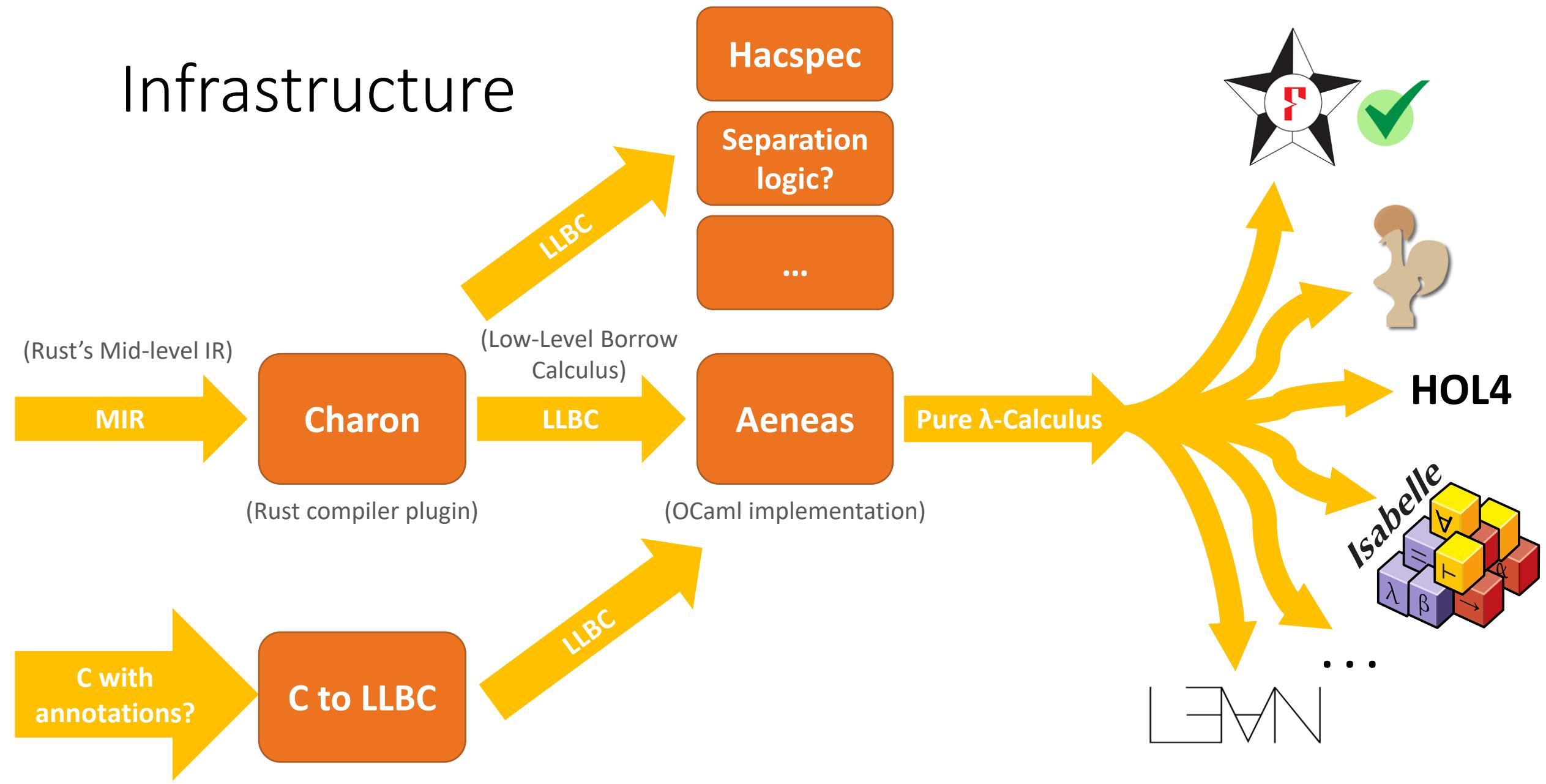
# Infrastructure



# Infrastructure



# Infrastructure



# Next steps

- Add at least one truly interactive backend (Coq, HOL4 , Lean)
- Extend supported subset (loops, traits...)
- Soundness proofs
- Ramping up verification:
  - B-epsilon-tree: ongoing
  - WireGuard VPN? MLS?

# Conclusion: Aeneas, in a nutshell

- **New point in the design space**
  - functional, modular, type-based translation
  - no annotations in source program
  - translation to existing proof frameworks for *extrinsic proofs*  
(now: F\*, ongoing : Coq, HOL4...)
- **Aim for verification in the large**
  - no “tour de force” on ninja unsafe code
  - rather, properties about large applications written in safe Rust
- **A user-friendly semantics for borrows**
  - lazy semantics, semantical analysis
  - yields a functional, executable translation, hence a program proof pipeline

# Conclusion: Aeneas, in a nutshell

- **New point in the design space**
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  - translation to existing proof frameworks for *extrinsic proofs*  
(now: F\*, ongoing : Coq, HOL4...)
- **Aim for verification in the large**
  - no “tour de force” on ninja unsafe code
  - rather, properties about large applications written in safe Rust
- **A user-friendly semantics for borrows**
  - lazy semantics, semantical analysis
  - yields a functional, executable translation, hence a program proof pipeline

## Questions ?

# Separation Logic Interpretation

```
fn choose<'a>(b : bool, x : &'a mut i32, y : &'a mut i32) -> &'a mut i32;
```

```
forall b x y x_v y_v.

{ (x ~> x_v) * (y ~> y_v) }

choose(b, x, y)

{ fun z ->
  // Returned value
  (z ~> choose_fwd b x_v y_v) *

  // Magic wand
  (forall z_v.
    (z ~> z_v) --*
    (let (x_v', y_v') = choose_back b x_v y_v z_v in
      (x ~> x_v') * (y ~> y_v')
    )) }
```

# Pure Translation – Lists & Recursion

```
pub enum List<T> {
    Cons(T, Box<List<T>>),
    Nil,
}

fn nth<'a, T>(l: &'a mut List<T>, i: u32) -> &'a mut T {
    match l {
        List::Nil => { panic!() }
        List::Cons(x, tl) => {
            if i == 0 { return x; }
            else { return nth(tl, i - 1); } } } }
```

```
let rec nth_fwd (l: list u32) (i: u32) : result u32 =
    match l with
    | Nil -> Fail
    | Cons x tl ->
        if i = 0 then return x
        else nth_fwd tl (i-1)

let rec nth_back (l: list u32) (i: u32) (n_x: u32) =
    match l with
    | Nil -> Fail
    | Cons x tl ->
        // If we reach the ith element:
        // replace "x" with its new value
        if i = 0 then Success (Cons n_x tl)
        // Otherwise continue diving into the list
        else
            match nth_back tl (i-1) n_x with
            | Fail -> Fail
            | Success n_tl ->
                // Replace "tl" with its new value
                Success (Cons x n_tl)
```

# Ongoing work: Loops / Translating blocks

There is a general problem of translating blocks of code in isolation:

```
// Rust
let x = 0;
if b { x = 1; }
else { x = 2; }
let y = x + 1;
...
```

```
// Current translation
let x = 0;
if b then
    let x = 1 in
    let y = x + 1 in
    ...
else
    let x = 2 in
    let y = x + 1 in
    ...
```

```
// Target translation
let x = 0;
let x = if b then 1 else 2 in
let y = x + 1 in
...
```

Idea: define a “join” operation on environments

# Discussion – Nested Borrows

- The real difficulty is dealing with **nested lifetimes**

```
fn id_mut_mut(ppx : &'a mut &'b mut u32)
-> &'a mut &'b mut u32 {
    x
}

let mut x = 0;
let mut px = &mut x;
let ppx = &mut px;
let mppx = &mut (*ppx);
let ppy = id_mut_mut(move mppx);
**ppy = 1; // Updates x
let mut y = 2;
*ppy = &mut y; // ← Now, (**ppx) == y
assert!(**ppx == 2); // This ends 'a
**ppx = 3; // ← We update y!
assert!(x == 1); // This ends 'b
assert!(y == 3);
```

```
// If we allow arbitrary type instantiations, we
// can turn any innocent looking function into
// something very complex:
fn id<T>(x : T) -> T { x }

// This is similar to what is on the left:
...
let mppx = &mut (*ppx);
let ppy = id<&mut &mut u32>(move mppx);
...
```

# Discussion – Nested Borrows – Use Cases

- What are the use cases for nested lifetimes (real question)?
- The only use case we are aware of is degenerate:

```
// This function is very idiomatic
fn f<'b, 'a>(ctx : &'b mut Ctx<'a>, ...) {
    // However the following “borrow overwrites” seem non-idiomatic:
    ctx.x = &mut ...;
}

fn g<'b, 'a>(...) -> &'b mut Ctx<'a>;
// The following “borrow overwrites” seem even less idiomatic:
ctx.x = &mut ...;
```

- Choose a subset which forbids the difficult operations?
- We are interested in use cases for “borrow overwrites”!

# Pure Translation – Proof of Concept (loops) (i)

```
fn list_iterator_has_next<'a, T>(l: &'a List<T>) -> bool {
    match l {
        List::Nil => {
            return false;
        }
        List::Cons(_, _) => {
            return true;
        }
    }
}

fn list_iterator_get_next<'a, T>(l: &'a mut List<T>) -> (&'a mut T, &'a mut List<T>) {
    match l {
        List::Nil => {
            panic!();
        }
        List::Cons(x, tl) => {
            return (x, tl);
        }
    }
}
```

```
let list_iterator_has_next (#a : Type0) (l : list a) : bool =
    match l with
    | [] -> false
    | x :: l -> true

let list_iterator_get_next_fwd (#a : Type0) (l : list a) : result (a & list a) =
    match l with
    | [] -> Fail
    | x :: tl -> Success (x, tl)

let list_iterator_get_next_back
    (#a : Type0) (l : list a) // Old value of the list
    (n_x : a) (n_tl : list a) :
    // return: the reconstructed list
    list a =
    n_x :: n_tl
```

# Pure Translation – Proof of Concept (loops) (ii)

```
fn list_iterator_test1(mut l: &mut List<u32>) {
    while list_iterator_has_next(l) {
        let (x, tl) = list_iterator_get_next(l); // Reborrow (i)
        *x = *x + 1;
        l = tl; // Reborrow (ii)
    }
}
```

```
let rec list_iterator_test1 (l0 : list u32) : result (list u32) =
    // > while list_iterator_has_next(l) {
    if list_iterator_has_next l0 then
        begin
            // > let (x, tl) = list_iterator_get_next(l);
            // This introduces a borrow on l, that we designate
            // as the borrow (i)
            tmp <-> list_iterator_get_next_fwd l0;
            let (x0, tl0) = tmp in

                x1 <-> check_add x0 1; // > *x = *x + 1;
                let l1 = tl0 in // > l = tl; // Borrow (ii)

                // Recursive call: end of the loop
                // We get the new value of l
                l2 <-> list_iterator_test1 l1;

                // Reborrows expire: propagate the values back
                let tl1 = l2 in // End borrow (ii)
                let l3 = list_iterator_get_next_back l0 x1 tl1 in // End borrow (i)
                return l3
            end
        else
            return l0
    end
}
```