

Formal Study of the French Tax Code's Implementation

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1 Introduction

2 SMT-based legislation study prototype

3 The M language

4 Conclusion

Code Général des Impôts, Article 197, I, 3, b, 3°

Le taux de la réduction prévue au premier alinéa du présent b est de 20 %. Toutefois, pour [...], le taux de la réduction d'impôt est égal à 20 % multiplié par le rapport entre :

- au numérateur, la différence entre 20 500 €, pour [...], ou 41 000 €, pour [...], et le montant des revenus mentionnés au troisième alinéa du présent b, et ;
- au dénominateur, 2 000 €, pour [...], ou 4 000 €, pour [...].

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When does the law specify an algorithm?

- Decision without human intervention
- No ambiguity
- Quantitative data (income, number of children, etc.)

An example of algorithmic translation

```
RATE = if (not [...1...]) then  
    20 %
```

“Le taux de la réduction prévue au premier alinéa du présent b est de 20 %. Toutefois, pour [...(1)...],...”

An example of algorithmic translation

```
RATE = if (not [...1...]) then
  20 %
else
  20 % * (
    (
      )
    /
    (
      )
  )
```

“ il est égal à 20 % multiplié par le rapport entre :

- au numérateur, ...;
- au dénominateur, ...”

An example of algorithmic translation

```
RATE = if (not [...1...]) then
  20 %
else
  20 % * (
    (
      (if [...2,3...] then 20 500 € else 41000 €)
      - INCOME
    )
    /
    (
      )
  )
)
```

“au numérateur, la différence entre 20 500 €, pour [...(2)...], ou 41 000 €, pour [...(3)...], et le montant des revenus mentionnés au troisième alinéa du présent b”

An example of algorithmic translation

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  20 %
else
  20 % * (
    (
      (if [...2,3...] then 20 500 € else 41000 €)
      - INCOME
    )
    /
    (if [...4,5...] then 2000 € else 4000 €)
  )
```

“au dénominateur, 2 000 €, pour [...(4)...], ou 4 000 €, pour [...(5)...].”

Formal study of the tax computation

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f should be studied formally!

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- Is f *increasing* with income?
- Is f *decreasing* with the number of children?
- What is f 's *derivative*? \Rightarrow marginal tax rate

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Case study: marginal tax rate

Model

Characteristic	Household before	Household after
Yearly income	R_0	$R_0 + \Delta_R$

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Effective withholding marginal rate

$$R_{\text{eff}} = \frac{\Delta_T + \Delta_B}{\Delta_R}$$

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Effective withholding marginal rate

$$R_{\text{eff}} = \frac{\Delta_T + \Delta_B}{\Delta_R}$$

Is there a household such that $R_{\text{eff}} \geq 70\%$?

Our hammer : SMT solving

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⇒ inefficient...
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Satisfiability modulo theories

Encoding of the tax computation within bitvector non-linear arithmetic and first-order logic.

⇒ Implemented in a prototype considering a simplified household model

Counter-example

The SMT solver finds something!

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Cohabiting couple with two high-schoolers (15 and 17-years-old), depending on second parent (jobless). Lives in zone II, monthly rent 897,75 €. Monthly raise of 250 €.

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Characteristic	Value before	Value after	Variation
Yearly income R	33 129,12 €	36 129,12 €	+ 3 000,00 €
IR	3 147,00 €	3 957,00 €	+ 810,00 €
PA	1 320,00 €	0,00 €	- 1 320,00 €
AF	1 584,00 €	1 584,00 €	0,00 €
ARS	806,00 €	0,00 €	- 806,00 €
BC	0,00 €	0,00 €	0,00 €
BL	0,00 €	0,00 €	0,00 €
APL	0,00 €	0,00 €	0,00 €
Take-home income N	33 692,12 €	33 756,12 €	+ 64,00 €
Marginal rate			97,9 %

More questions...

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`https://gitlab.inria.fr/verifisc/verifisc-python`

SMT-verified fiscal legislation?

Avantages

- Systematic analysis
- Very expressive model
- No data needed

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Scalability challenge

- Complexity \Rightarrow resources and compute time
- SMT queries need optimizing
- Going beyond prototyping

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From prototype to production: leveraging DGFIP's work

The DGFIP¹ released a part of its tax computation system in 2016:

- <https://framagit.org/dgfip/ir-calcul>
- Written in M, custom language
- Computes income taxes for private individuals
- No official, publicly available documentation (no grammar, no semantics, ...)
- 2017 version: 48 files, 92,000 lines of code

¹“Direction Générale des Finances Publiques”: administration in charge of the tax system.

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Let's use the DGFIP's work!

- Kept up to date with new legislation by DGFIP
- "Official code"

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A tax-specific DSL

Values in M

- Booleans or undef
- Floating-point numbers or undef
- Fixed-size arrays of homogeneous values

A tax-specific DSL

Values in M

- Booleans or `undef`
- Floating-point numbers or `undef`
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Essence of M

- Statements: mostly variables assignments; exceptions can be raised
- Expressions: thresholds (`min`, `max`) and arithmetic computations
- A *lot* of values $v \in \{0, 1\}$ used as booleans
- A *few* loops, all having a known, constant iteration number

A tax-specific DSL: an example

Variable declaration

```
IRNETBIS : calculee primrest = 0 : "IRNET avant bidouille du 8ZI";
```

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Computations

```
regle 221220:  
application : iliad ;  
IRNETBIS = max(0, IRNETTER -  
               PIR * positif(SEUIL_12 - IRNETTER + PIR)  
               * positif(SEUIL_12 - PIR)  
               * positif_ou_nul(IRNETTER - SEUIL_12));
```


A formal semantics for M

- Reverse-engineering of the semantics
- μ M semantics written in Coq

Looks quite simple:

$$\frac{\text{D-VAR} \quad \Omega(x) = v}{\Omega \vdash x \Downarrow v}$$

$$\frac{\text{D-COND-TRUE} \quad \Omega \vdash e_1 \Downarrow \text{true} \quad \Omega \vdash e_2 \Downarrow v_2}{\Omega \vdash \text{if } e_1 \text{ then } e_2 \text{ else } e_3 \Downarrow v_2}$$

$$\frac{\text{D-INDEX} \quad \Omega(x) = (v_0, \dots, v_{n-1}) \quad \Omega \vdash e \Downarrow r \quad r \in \llbracket 0, n-1 \rrbracket}{\Omega \vdash x[e] \Downarrow v_r}$$

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Except for the undef value...

A formal semantics for M

The undefined value

Used for:

- default inputs
- runtime errors
- missing cases in inline conditionals

undef-related rules:

$f_1 \odot f_2, \odot \in \{+, -, \times\}$		undef	$f_2 \in \mathbb{F}$
undef		undef	$0 \odot f_2$
$f_1 \in \mathbb{F}$		$f_1 \odot 0$	$f_1 \odot f_2$

A formal semantics for M

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$f_1 \odot f_2, \odot \in \{+, -, \times\}$	undef	$f_2 \in \mathbb{F}$	$f_1 \div f_2$	undef or 0	$f_2 \in \mathbb{F}, f_2 \neq 0$
undef	undef	$0 \odot f_2$	undef	undef	0
$f_1 \in \mathbb{F}$	$f_1 \odot 0$	$f_1 \odot f_2$	f_1	undef	$f_1 \div f_2$

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undef	undef	$0 \odot f_2$	undef	undef	0
$f_1 \in \mathbb{F}$	$f_1 \odot 0$	$f_1 \odot f_2$	f_1	undef	$f_1 \div f_2$

D-COND-UNDEF

$\Omega \vdash e_1 \Downarrow \text{undef}$

$\Omega \vdash \text{if } e_1 \text{ then } e_2 \text{ else } e_3 \Downarrow \text{undef}$

The Mlang compiler

- Implementation following the semantics
- Input: DGFIP's codebase and query file
- Usecases:
 - Tax computation given input parameters (given in the query file)
 - Compilation to Python, Java; query specifying extraction assumptions
- Written in OCaml
- Freely available: <https://gitlab.inria.fr/verifisc/mlang>

Code optimization

Classical compiler optimizations

- Global value numbering
- Dead code elimination
- Partial evaluation

Simplified simulator case

A case², with less inputs: 2460 to 231 variables

Compiler optimizations: 46,364 variables down to 1,600

²https://www3.impots.gouv.fr/simulateur/calcul_impot/2019/simplifie/index.htm

Experimental evaluation

The DGFIP provided us with around 500 tests³:

- Each test provides inputs and assertions on some variables (median LOC: 597)
- We pass a fifth of the tests (558 to 674 LOC)
- Failed tests: we are missing some constants not provided by the DGFIP
- Awaiting technical answers from the DGFIP

³They should eventually be public.

Future work

- Pass all tests
- Translation to Z3 (need to take undefs into account)
- Abstract interpretation for further optimization (Z3 and other)

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Related work

Academic research

- As old as 1956 [1]!
- Prolog for determining British nationality [7]
- US tax code formalization using default logic [4, 5]
- Smart contracts [3]

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Private ventures

- Formal vindications [2]
- Imandra [6]
- Legalese

Conclusion

Current work and beyond:

- SMT-based tax impact study
- Generating verified code to compute taxes, get it to production
- Upstream formalization to legislative text production

Open-source code: <https://gitlab.inria.fr/verifisc>

JFLA submission: <https://hal.inria.fr/hal-02320347>

Thanks to DGFIP and in particular Christophe Gaie and his team for collaborating!

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